Security in the Ambient Computational Environment

S.Vidyaraman

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Committee: Dr.Joseph. B Evans (Chair) Dr.Gary.J Minden Dr.Arvin Agah

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Acknowledgements

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Overview

- Background
- Security Issues Addressed
- Security Services Implemented
- Typical Scenarios
- Analysis
- Summary & Future Work
- Q&A



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Background

- ACE : Ambient Computational Environment
- Its all about reinventing the 4 wheels of the car. But then
- Entities in ACE
 - ACE Services
 - ACE Users



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- Services communicate within themselves.
 - Network Commands
 - Data Streams (Audio and Video)
- Users
 - Authentication
- The Users Workspace is a VNC Session.
- How do we identify both Users and Services?



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Security Services Implemented

• Remote Connection Manager



- Certificate Authority
- Certificate Distribution System
- Key Manager



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Security Services Implemented

- <u>Remote Connection Manager</u>
 - Functionality
 - DH Key Establishment
 - <u>SPEKE Protocol</u>
- Certificate Authority
- Certificate Distribution System
- Key Manager



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Remote Connection Manager

- Gateway to the ACE Domain from Outside
- Functions:
 - Authenticate the user
 - Establish a shared session key



- At present, it implements the SPEKE protocol
 - A Variant of the Diffie-Hellman Key establishment
 - One of the *strong* authentication mechanisms with (even) weak passwords
 - Minimum (3) number of passes
 - Protects against dictionary attacks



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SPEKE Protocol

Secure		Alice		Bob
Password-authenticated		All operations are mod p		
Exponential		$\mathbf{Q}_{\mathbf{A}} = \mathbf{S}^{(2 \mathbf{R}} \mathbf{A})$	\rightarrow	
Key	Кеу		÷	$\mathbf{Q}_{\mathbf{B}} = \mathbf{S}^{(2 \mathbf{R} \mathbf{B})}$
Exchange	Exchange	$\mathbf{K} = \mathbf{Q}_{\mathbf{B}}^{(2 \mathbf{R} \mathbf{A})}$		$\mathbf{K} = \mathbf{Q}_{\mathbf{A}}^{(2 \mathbf{R} \mathbf{B})}$
8-		Abort if K< 2		Abort if K< 2
The generator g is now the				
squared hash of the password S			\	$\mathbf{V}_1 = \mathbf{h}(\mathbf{h}(\mathbf{K}))$
	Verification	$\mathbf{V}_2 = \mathbf{h}(\mathbf{K})$	\rightarrow	
		Abort if V ₁ != h(h(K))		Abort if V ₂ != h(K)
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Security Services Implemented

• Remote Connection Manager



- Public Key Infrastructure (PKI based Services)
 - <u>Certificate Authority</u>
 - <u>Certificate Distribution System</u>
- Key Manager



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Certificate Authority



- Provides identification to users and daemons
- Issues X509 digital certificates to users & daemons
- Revokes the user / daemon certificate when necessary
 - Creates a CRL for all the certificates revoked
- Notifies the issued & Revoked Certificates to the Certificate Distribution Daemon



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Certificate Distribution System



- Function: To distribute all valid user / daemon certificates
- Answers queries from ACE services regarding validity of certificates
- Publishes the list of valid certificates and the Certificate Revocation List (CRL) on a publicly accessible LDAP service



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ACE Root Certificate

- Same Issuer and Subject
- Essentially a self signed Certificate
- Signature Algorithm: md5withRSA
- Thumbprint Algorithm: sha1

ihow: <all></all>	
Field	Value
🛅 Serial Number	01
🚍 Signature Algorithm	md5RSA
E Issuer	Research & Developement, ACE: ITTC Doma
Calid From	Monday, April 08, 2002 4:32:26 AM
🗖 Valid To	Friday, June 07, 2002 4:32:26 AM
Subject	Research & Developement, ACE: ITTC Doma
Public Key	RSA (2048 Bits)
<u> </u>	
CN = ACE: ITTC Domain T = Certificate Authority C = USA L = Lawrence E = ace@ittc.ku.edu 0 = University of Kansas	
	Edit Properties



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ACE Certificate Revocation List

Certificate Rev	ocation List Information	<u>R</u> evoked certificates:	
8-1		Serial Number	Revocation Date
Field	Value	03	Wednesday, May 08, 20
	V2		
- Issuer	Research & Developement, ACE: I		
Effective Date	Wednesday, May 08, 2002 4:34:22		
💳 Next Update	Tuesday, May 07, 2002 4:32:30 AM		
💳 Signature Algorithm	md5RSA	Revocation entry	
		Field	Value
		Serial Number	03
		Revocation Date	Wednesday, May 08, 2002 4:34:22 AM
/alue:			
JU = Hesearch & Develop CN = ACE: ITTC Domain	pement		
F = Certificate Authority		Wednesday, May 08, 2	UU2 4:34:22 AM
C = USA			
. = Lawrence F = ace@itto ku edu			
) = University of Kansas			
	OK		OK

Security Services Implemented

• Remote Connection Manager



- Certificate Authority
- Certificate Distribution System
- <u>Key Manager</u>



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- One time session
- Conference
- All issued keys are stored in a PBE encrypted keystore



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Typical Scenarios

• Remote Authentication

• Certification Process



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Possible Remote Authentication Procedure



Remote Authentication Process







- Password / IButton ID / Fingerprint ID
- Standard encryption



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Analysis



- 3. What new problems have been added?
 - Addition overhead of managing a limited PKI
 - Additional vulnerability to social engineering problems
 - Passwords can be changed once a compromise is detected
 - Not true with IButton and Fingerprint data
- Extraneous issues!
 - Java
 - API calls & Key lengths





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Summary

- The following services have been prototyped in this thesis
 - A Rudimentary Key Manager
 - A Certificate Authority
 - A Certificate Distribution System
 - A Remote Connection Manager
- But thenSecurity is a process, not a product.



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Future Work

- Implement a (m,l)- threshold b-secure t-group key distribution scheme
 - Number of centers: *m*
 - Minimum number of centers required: *l*
 - (*l*-1) center & b user compromise doesn't compromise the system
- Better storage system for CA Keys and Certificates



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X.509 Digital Certificate Fields

Certificate field	Description			
Version	The X.509 version number.			
Serial number	The unique serial number that the issuing certification authority assigns to the certificate. The serial number is unique for all certificates issued by a given certification authority.			
Signature algorithm	The hash algorithm that the certification authority uses to digitally sign the certificate.			
Issuer	Information regarding the certification authority that issued the certificate.			
Subject	The name of the individual or certification authority to which (whom) the certificate is issued. This may be a full name and e-mail name or some other personal identifier.			
Public key	The public key type and length associated with the certificate.			
Thumbprint algorithm	The hash algorithm that generates a digest of data (or thumbprint) for digital signatures			
Thumbprint	The digest (or thumbprint) of the certificate data.			
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SPEKE Vs DH-EKE

modulus p is hugediscrete log attackD Stest Qx != 0, when un-encryptedforcing K=0D Sp-1 has large prime factor qPohlig-Hellman log computationD Sencrypted Qx randomly padded.leakage from Es(Qx)Dbase is primitive root of ppartition attack on Es(Qx)Dbase is a generator of qpartition attack on QxSbase = Sx mod ppassword-in-exponent attackSfirst receiver of verification of K mustsfinding password S using chosen Rx, Qx, and password dictionaryDuse one-way hash of Knarrowing attacks on Ek(Qx)Dhigh bits of p must be 1partition attack on Ek(Qx)DReceiver of clear Qx abort if K is small order. or Encrypt QA, QB.subgroup confinement of KSAbort if K has small ordersubgroup confinement of KS	Constraint	Prevents Attack by:	Applies to
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	Abort if K has small order	subgroup confinement of K	S



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Navigation

- ACE Entities
 - <u>Services</u>
 - <u>Users</u>
- Services
 - <u>Remote Connection Manager</u>
 - <u>Certificate Authority</u>
 - <u>Certificate</u>
 - <u>Certificate Revocation</u>
 <u>List</u>
 - <u>Certificate Distribution</u> <u>System</u>
 - <u>Key Manager</u>



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- <u>Diffie-Hellman</u>
- Scenarios
 - <u>Remote Authentication</u>
 - <u>Certification</u>
- <u>Analysis</u>

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• <u>Distributed Key</u> <u>Manager</u>