Scalable Emulation of IP networks through Virtualization

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Information and Telecommunication Technology Center

Talk Content

- Introduction
- Virtual Network Framework
- Design of Virtual Network Elements
- Evaluation of Virtual networks
- Summary



Introduction – Problem

- Ubiquitous IP networks Data, Video & Voice
- Need to study and test new protocols
- Large scale networks
- Current testing methods:
 - Simulation
 - Physical testing



Introduction – Existing methods

- Simulation
 - ns2, OPNET
 - Side-effects of OS interactions ignored
 - Management complexity ignored
 - Change in focus
- Physical Testing
 - Equipment/Infrastructure costs
 - Erroneous extrapolations of small tests



Introduction - Goals

- Design & Implement Network Emulation Framework to solve current problems
- Test with realistic network loads
 - Generated by real utilities (e.g. ttcp, ftp, telnet, etc.)
 - Synthesized loads through Netspec
- Compare results with results from physical network



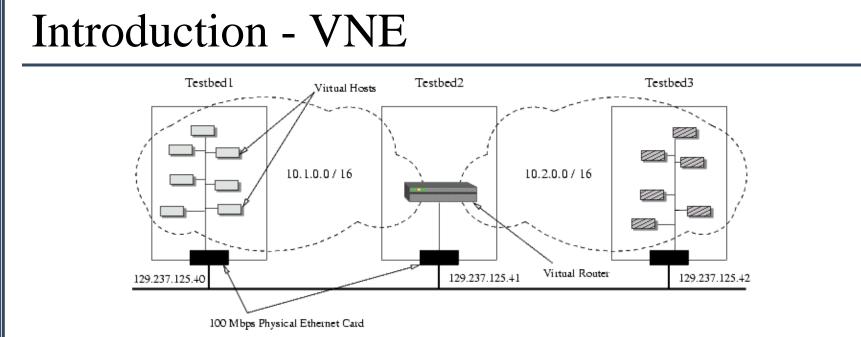
Introduction – Virtual Network Elements

• Definition:

A virtual network element (VNE) is a software object that emulates the functions of network elements such as hosts and routers.

- Modules inserted into the Linux network protocol stack transparently
- New layer added to protocol processing sequence: *Virtual Network layer* (VNL)





- Simple application of virtual network
- Virtual network traffic multiplexed over physical interface(s)
- VNL handles mux/demux



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Virtual Network Framework (VNF)

- Three basic elements of a network: host, router & link
- Host and Router emulated by virtual host and virtual router code in VNL
- Virtual link implemented using link throttling techniques of Linux traffic control
- Σ throughput(virtual elements) $\leq \Sigma$ throughput(physical interfaces)

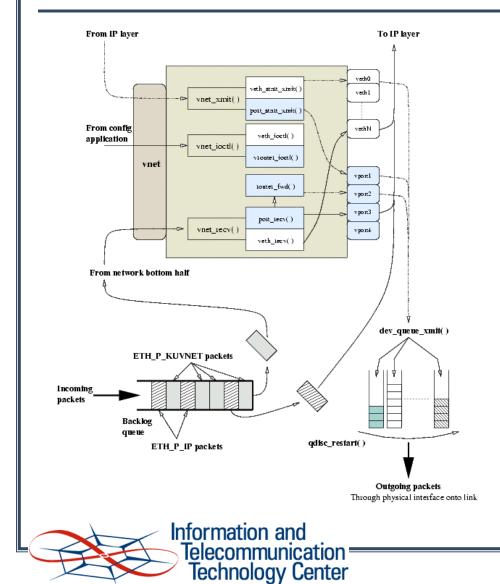


VNF: Design Considerations

- Socket-layer compatibility
- Creation/Deletion/Configuration
- Arbitrary mapping of virtual elements to physical hosts
- Virtual routing decisions
- Network emulation ability



VNF: Architecture



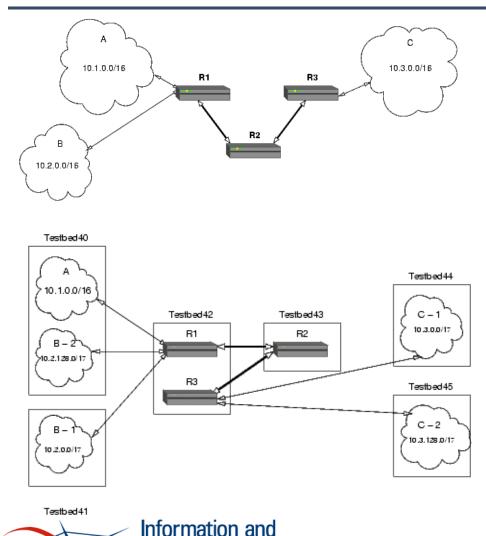
- Multiple virtual hosts and virtual routers share VNL
- ETH_P_KUVNET is the packet type
- Each of the virtual devices can have a queue attached

VNF: Capabilities

- Multi-homed (virtual) hosts
- Split subnets across physical machines
- Supports almost arbitrary mapping of virtual elements to physical hosts through *subnet maps*
- Subnet map identical to routing table
- VNL inserts a new header: VNET header between IP and MAC header



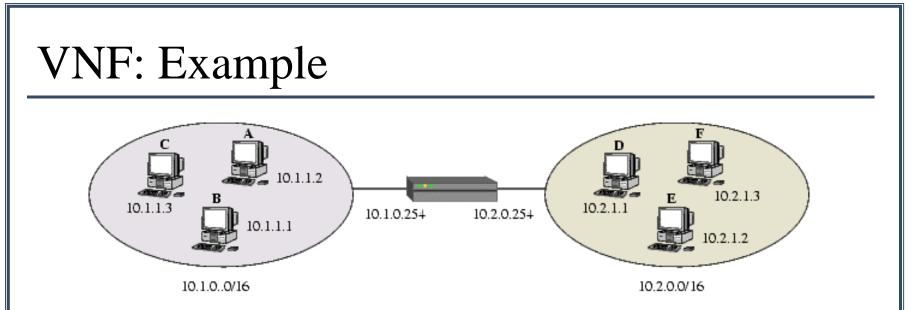
VNF: Capabilities: Split-subnet



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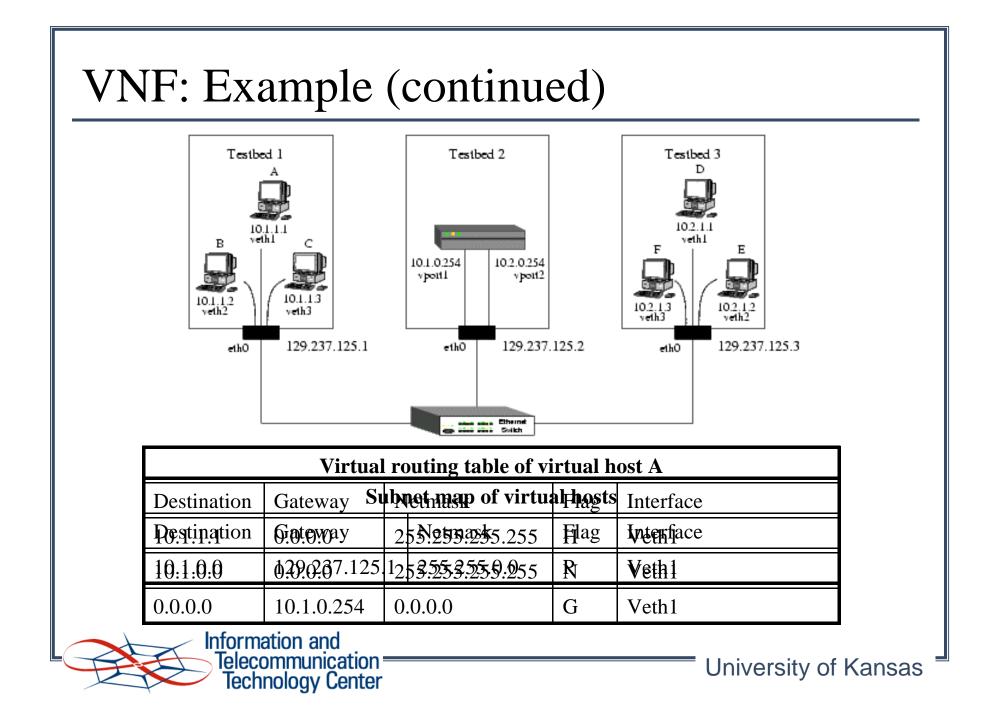
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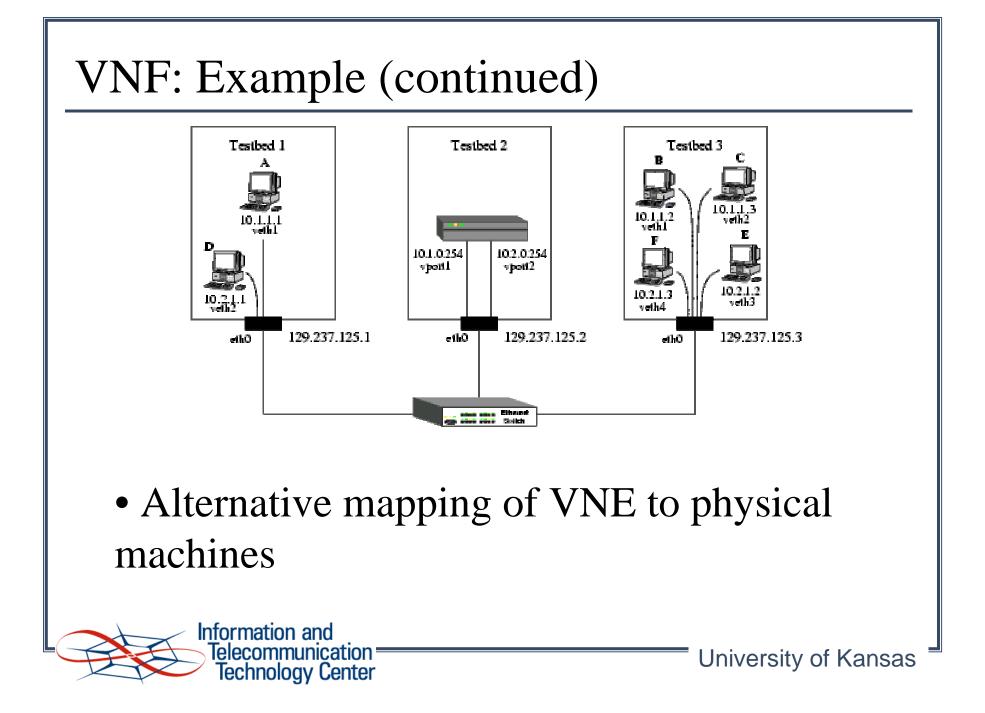
- Flexibility in placing VNEs on physical machines
- Load/Application /Characteristicbased mapping



- Depicts virtualization based on network application
- Depicts working of *split-subnet* mapping
- A, D: Servers
- $A \leftrightarrow (E, F), D \leftrightarrow (B, C)$







Design of Virtual Network Elements

- Implemented as Linux network device driver
- Configured through ioctl()s
- Netspec-based configuration
- Shows virtual interface statistics through tools such as ifconfig, ip, etc.
- Supports packet capture tools such as tcpdump



VNE design: Virtual Host

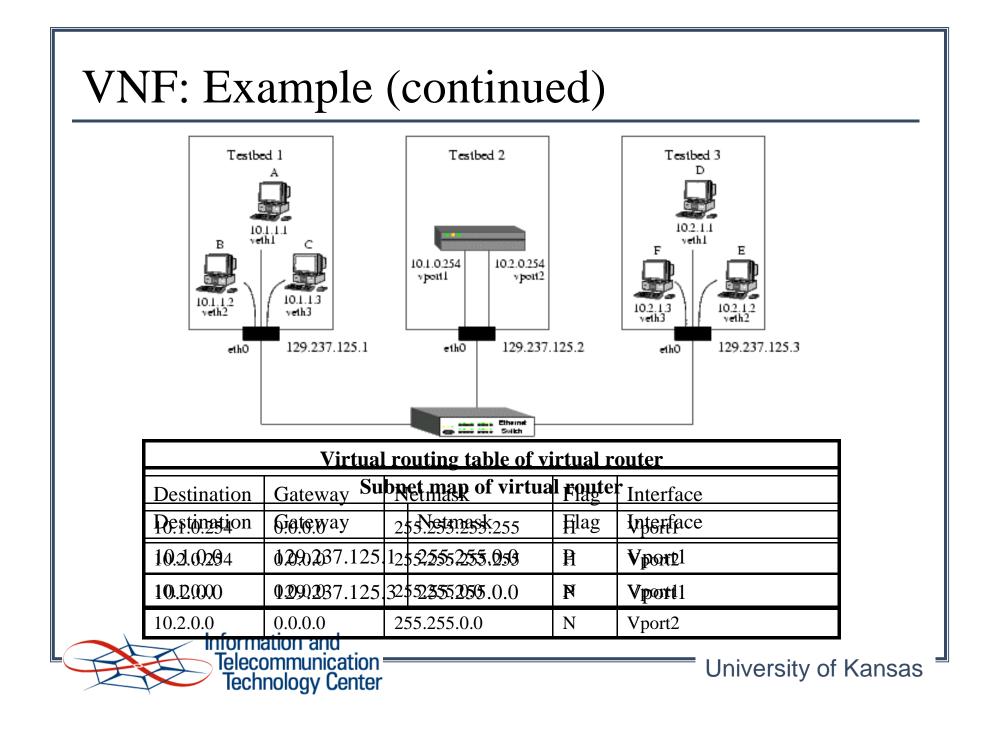
- Has an IP address
- Virtual routing table contains gateway entry
- Subnet map table contains location of virtual router emulating the gateway
- Acts as source or sink
- Socket applications bind to it



VNE: Virtual Router

- Each port has an IP address
- Virtual routing table contains entries to other routers or to subnets
- Subnet map table contains location of virtual router and subnets
- Only a gateway for the packets
- Socket applications typically don't bind to it (exception: RSVP daemon)





Evaluation of Virtual Networks

- Control plane of emulated network remains the same as physical network
 - Identical software
 - Identical signaling costs
- Need to confirm verity of data plane results
- Results of Physical tests vs. Emulation tests
- Diffserv and Intserv networks used for exercising VNF

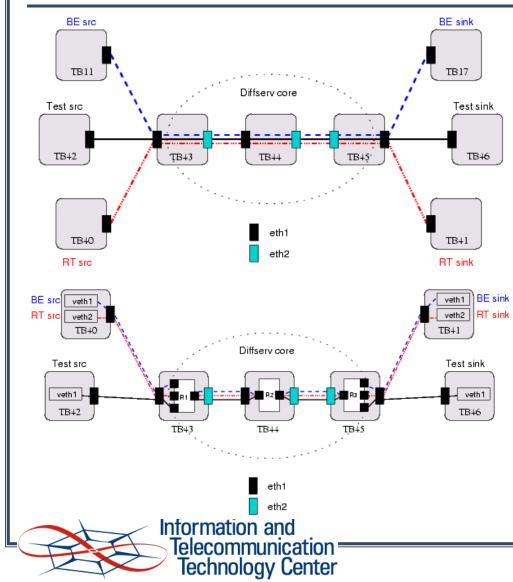


Evaluation of Virtual Networks

- Diffserv relies on Linux traffic control(*tc*), hence works with VNE with minor modifications
- Intserv relies on *tc* and RSVP signaling, RSVP required some porting to understand virtual routing



Diffserv - Network Topology (9 elements)



- Link bandwidth 100Mbps in access network & 10Mbps inside core
- Routers Emulated on high speed
 Pentium III, 1GHz,
 1GB RAM Linux
 systems

Diffserv - Network Topology (9 elements)

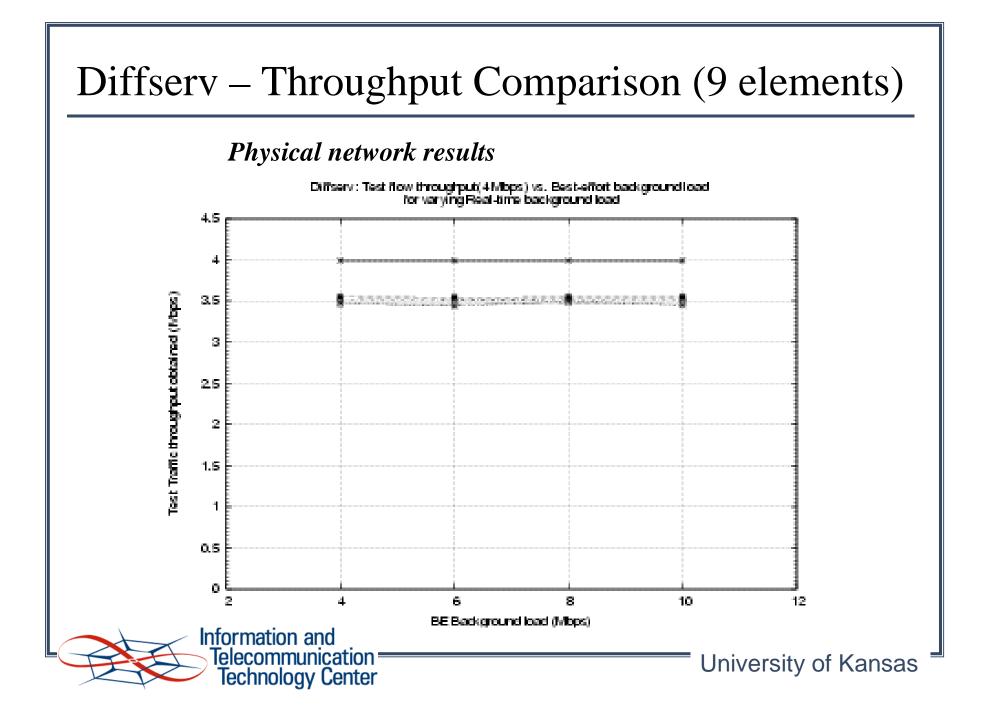
- Used to validate working of Diffserv
- Throughput of 'Test' stream measured in presence of background RT and BE load
- Netspec-generated CBR traffic using UTIME patches
- tcpdump output captured at source and sink, merged and diff'ed

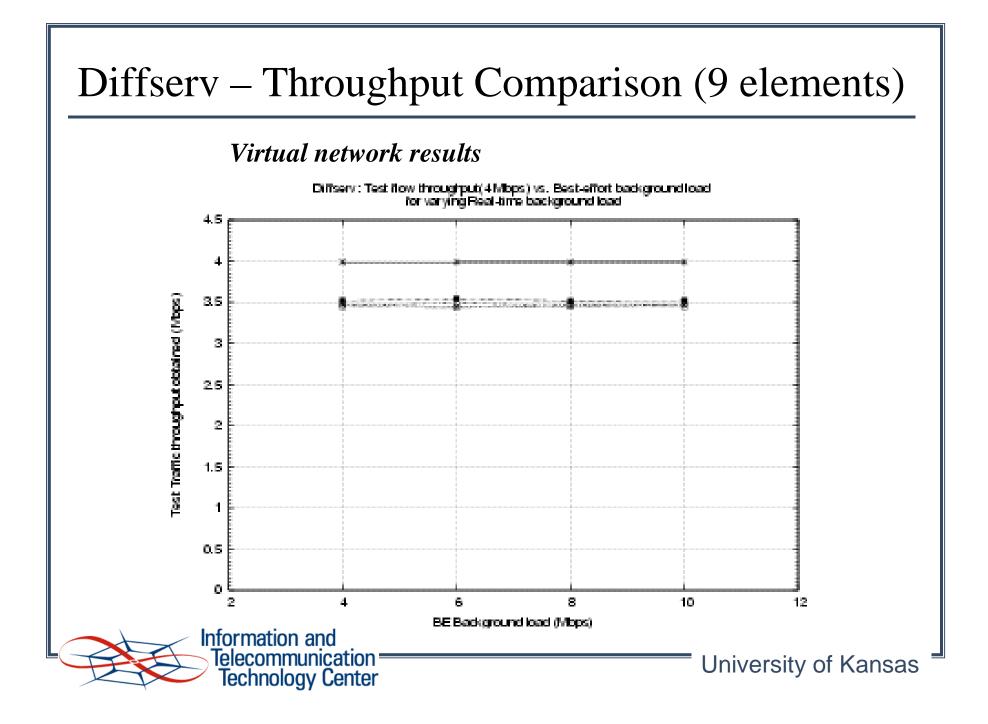


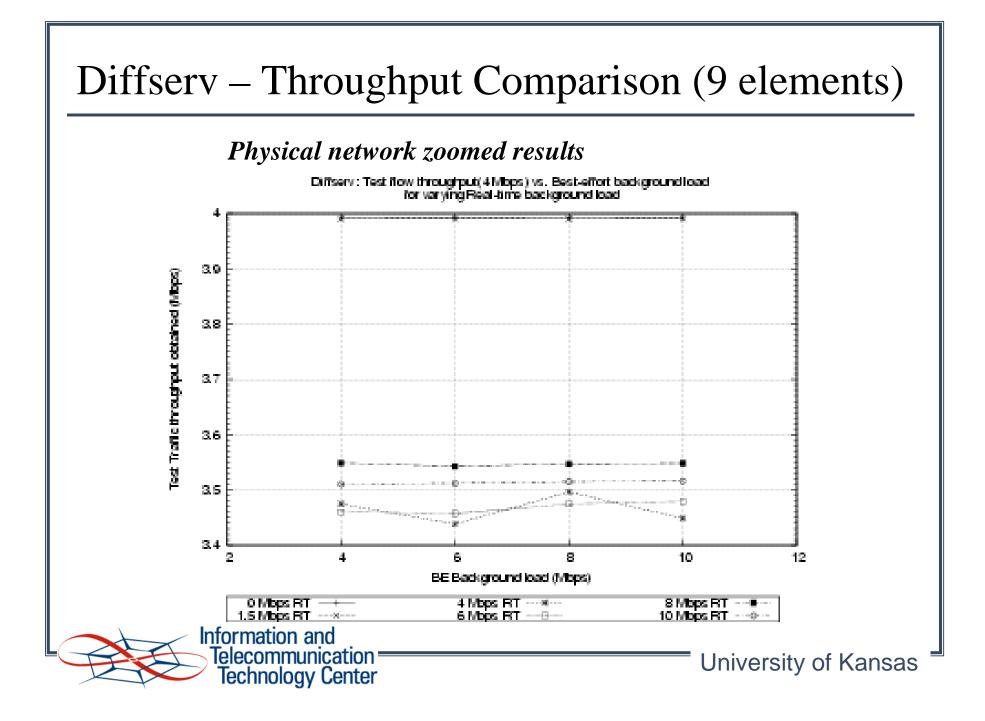
Diffserv - Network Test parameters

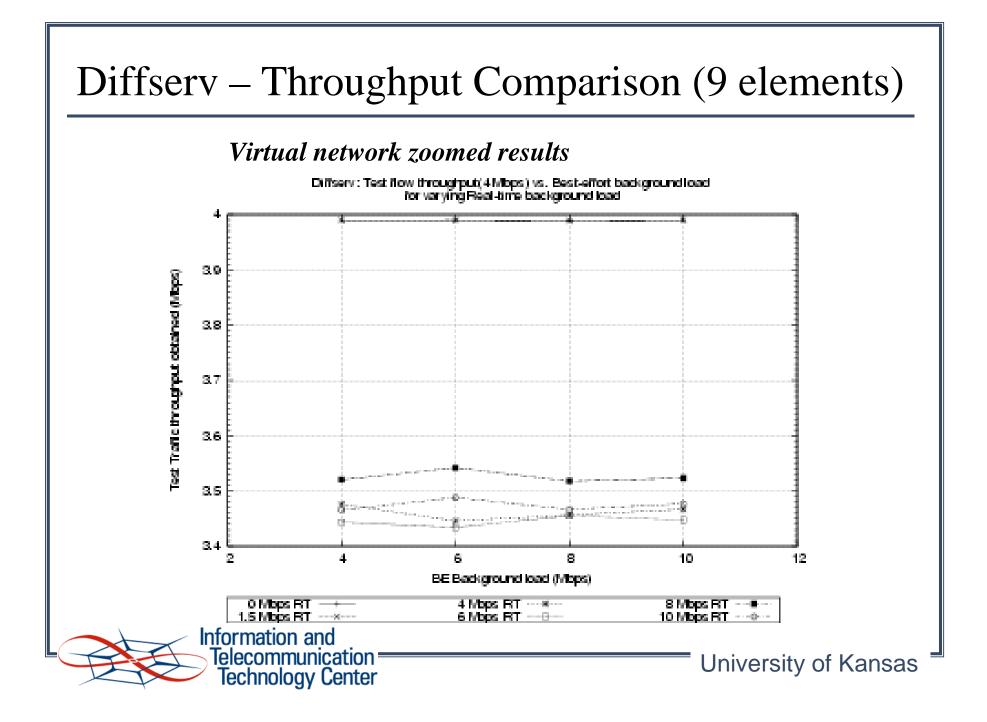
Data Plane	• Delay	
	• Throughput	
	• HTB queuing discipline	
(core routers)	• Best Effort class = 4Mbps	
Diffserv Parameters	• Real time AF class = 6Mbps	
	• Test CBR traffic = 4Mbps	
Traffic	• BG-RT traffic = 0-10Mbps	
	• BG-BE traffic = 4-10Mbps	



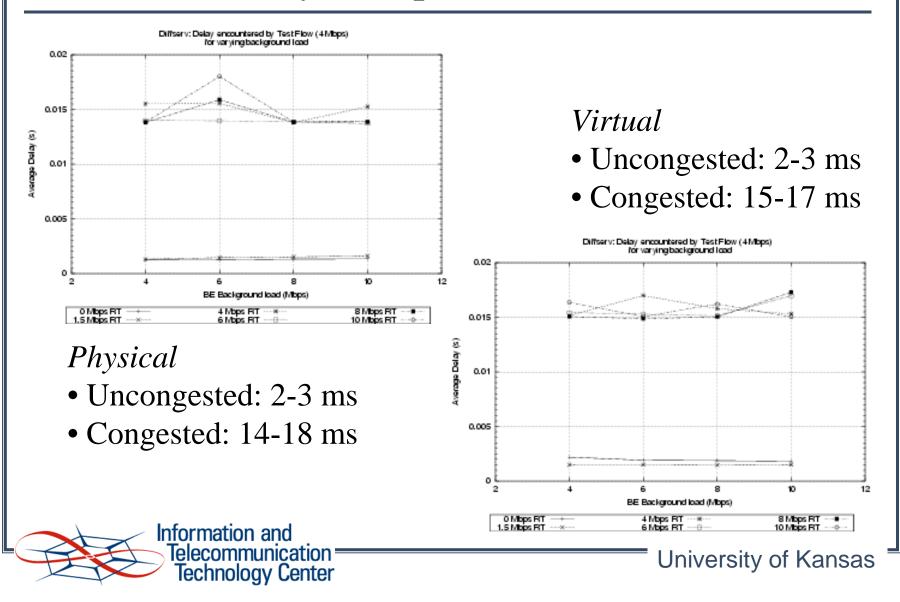


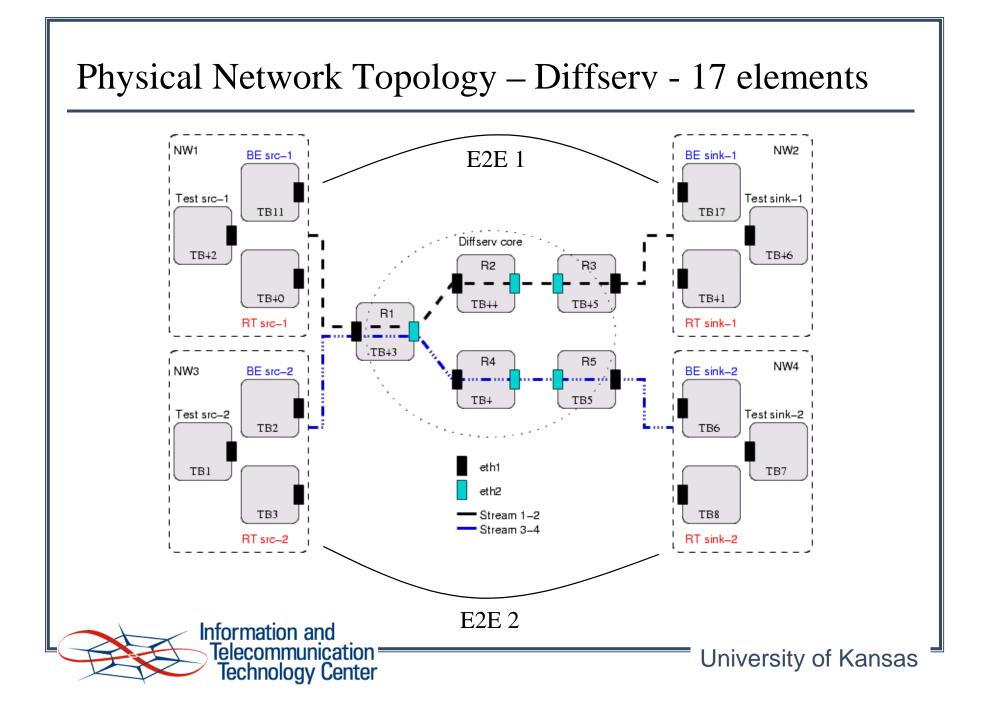






Diffserv – Delay Comparison (9 elements)

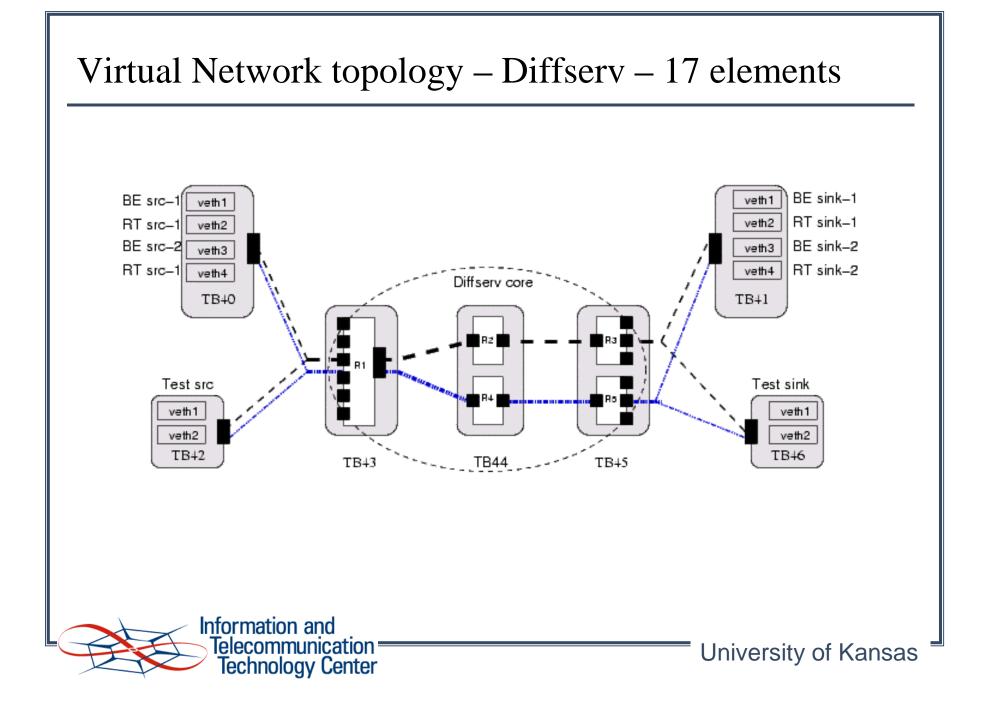




Virtual Network topology – Diffserv – 17 elements

- Physical testing not performed due to shortage of machines
- Ideal case for using VNF
- Need to compare Diffserv properties observed in Physical networks with those observed in Emulated networks





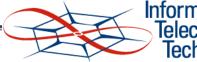
Virtual Network topology – Diffserv – 17 elements

- One greedy customer does not affect other customers of network
- Throughput of 2 'Test' streams measured in presence of background RT and BE load from their respective networks
- tcpdump output captured at source and sink, merged and diff'ed



Diffserv - Network Test parameters

Traffic	 BG-BE traffic = 2-6Mbps BG-RT traffic = 0-6Mbps Test CBR traffic = 4Mbps 	
Diffserv Parameters (core routers)	 2 Real time AF classes = 6Mbps each Best Effort class = 4Mbps HTB queuing discipline 6 Mb RT traffic is threshold (Test + BG) 	
Data Plane	ThroughputDelay	



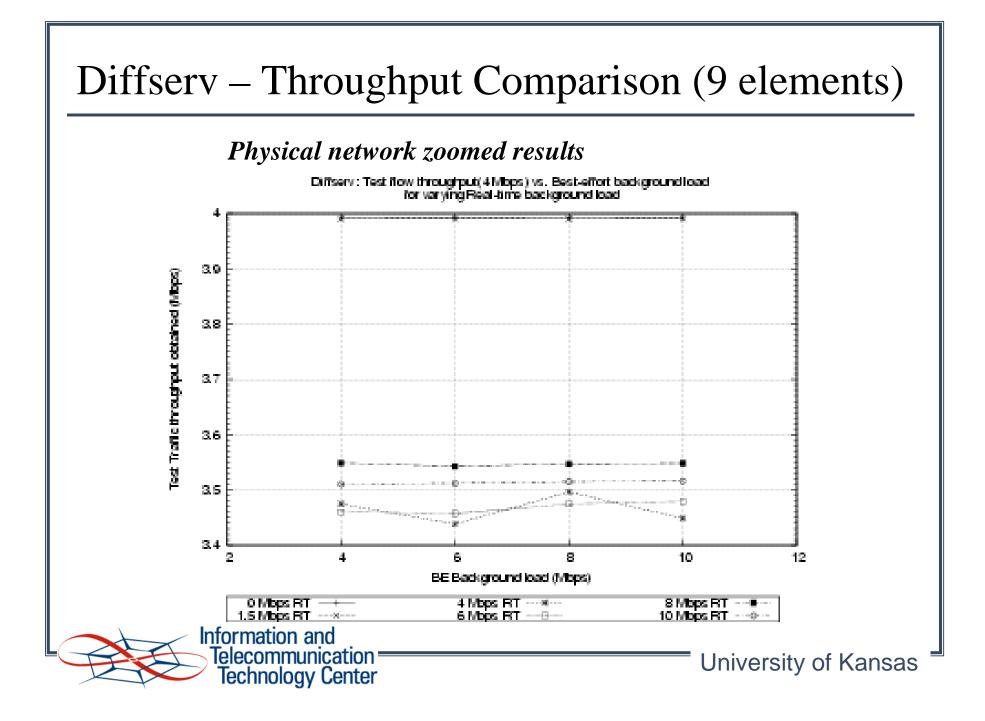
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Diffserv – throughput Results (17 elements)

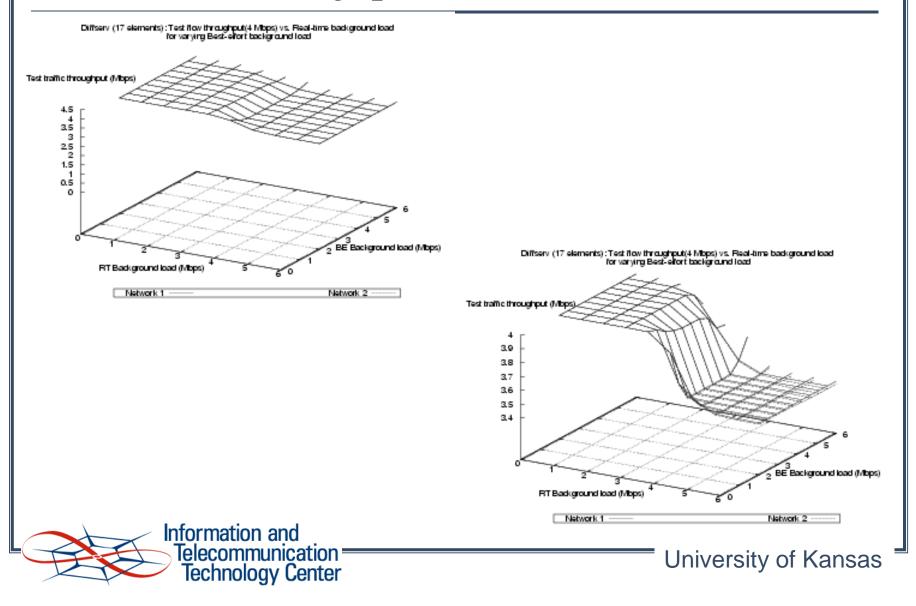
- #1-6 show equal traffic on both networks
- #7-8 show E2E #2 being greedy
- One greedy customer does not affect others in



#	E2E	BG-RT (Mbps)	BG-BE (Mbps)	throughput (Mbps)
1	1	0	2	3.9901
	2	0	2	3.9901
2	1	0	4	3.9904
	2	0	4	3.9901
3	1	1.5	2	3.9901
	2	1.5	2	3.9901
4	1	1.5	4	3.9900
	2	1.5	4	3.9897
5	1	4	4	3.4735
	2	4	4	3.4648
6	1	6	6	3.4752
	2	6	6	3.4551
7	1	1.5	2	3.9901
	2	4	4	3.4799
8	1	1.5	4	3.9901
	2	6	6	3.4591



Diffserv – Throughput Results (17 elements)



Diffserv Data Plane Evaluation – Results

- Diffserv behaves similarly in physical and emulated (virtual) networks
- Results of throughput/delay tests on emulated network similar to those of physical network
- Very minor changes to code to get Diffserv to work with VNEs



Emulating Intserv networks

- Emulated network is identical to Diffserv network
- Diffserv traffic classes replaced by RSVP daemon which does dynamic resource reservation
- RSVP daemon modified:
 - To understand virtual routing
 - To enable many instances to run on a physical machine bound to specific VNEs



Emulating Intserv networks

- Intserv network successfully emulated
- Results on Physical network not reproducible for multiple iterations of tests
 - RSVP daemon uses CBQ
 - Linux CBQ implementation tries *estimation* to schedule packets, does not give consistent results
 - HTB implementation for RSVP non-trivial
- Data plane could not be verified
- Demonstrates clean interface of VNF that allows complex applications to use it



Limitations of VNF

- Sum of throughputs of VNEs on a physical machine must be less than sum of throughput of all physical interfaces; overcome using virtual time techniques introduced by ProTEuS
- Doesn't allow 'connected' NEs to be emulated on same physical machine if packet needs to pass through queuing code; can be overcome by modifying queuing code



Summary

- VNF designed and implemented
- Tested with non-trivial IP networks such as Diffserv and Intserv networks
- Programming model allows easy 'porting' of applications to work with VNF
- Larger Diffserv networks successfully emulated
- Intserv networks emulated functionally, but data plane could not be verified



